

Zen and The Art of Grammar Maintenance

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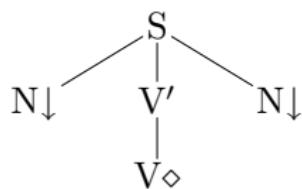
Outline

1 Motivations

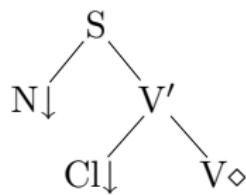
2 Metagrammars as Logic Programs

3 XMG Formalism

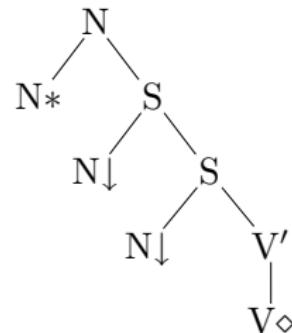
Factorization and Reuse



Jean mange une pomme

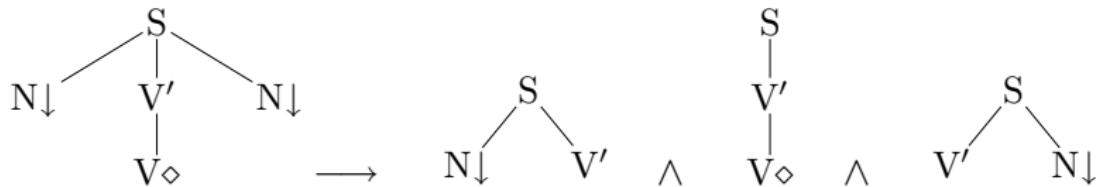


Jean la mange



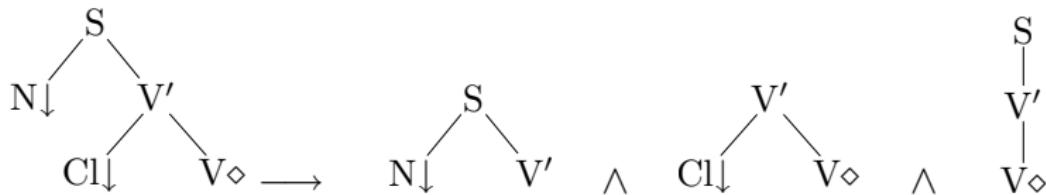
la pomme que Jean mange

Jean mange une pomme



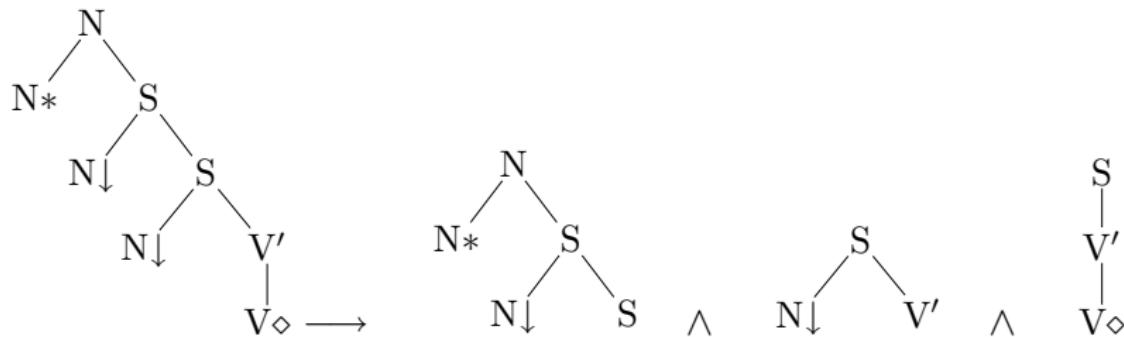
- Canonical Subject
- Active Verb Form
- Canonical Object

Jean la mange



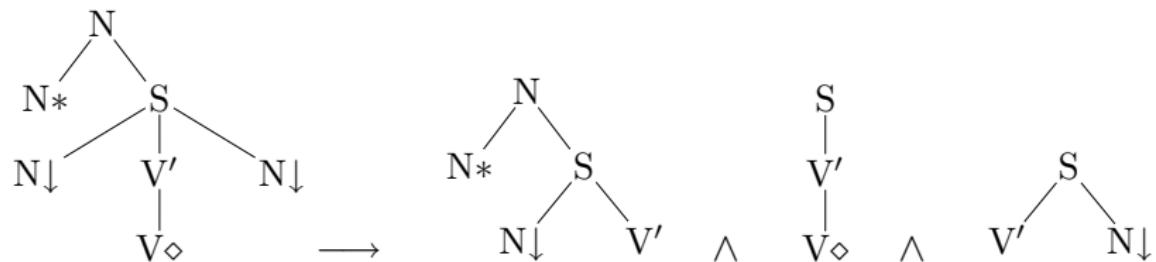
- Canonical Subject
- Clitic Object
- Active Verb Form

La pomme que Jean mange



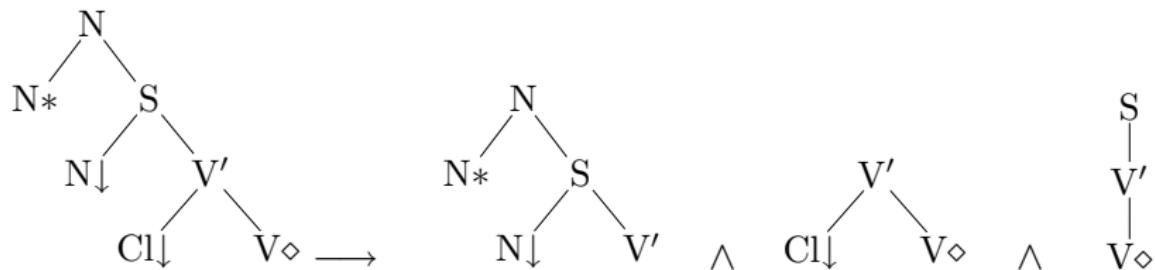
- Relative Object
- Canonical Subject
- Active Verb Form

Jean qui mange la pomme



- Relative Subject
- Active Verb Form
- Canonical Object

Jean qui la mange



- Relative Subject
- Clitic Object
- Active Verb Form

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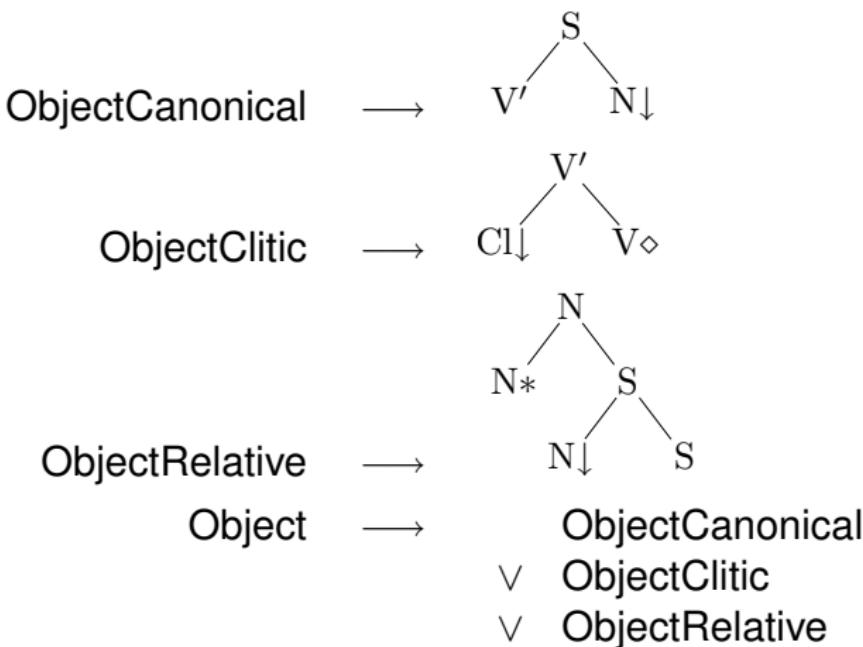
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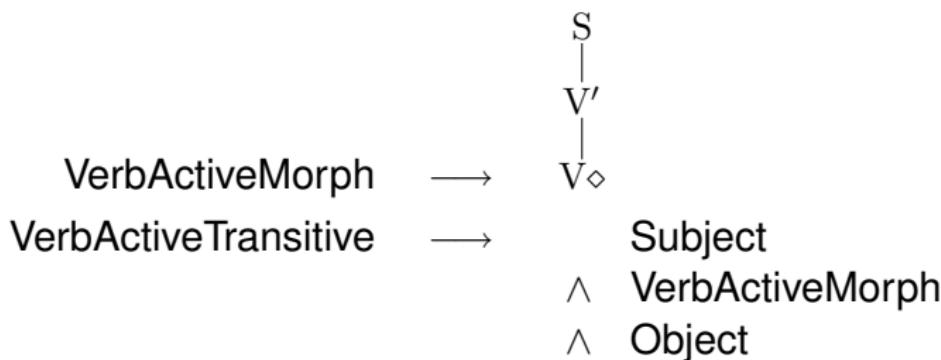
Subject

SubjectCanonical	\longrightarrow	<pre> graph TD S1[S] --> N1[N] S1 --> V1[V'] </pre>
SubjectRelative	\longrightarrow	<pre> graph TD N2[N] --> N3[N*] N2 --> S2[S] S2 --> N4[N] S2 --> V2[V'] </pre>
Subject	\longrightarrow	$\begin{array}{c} \text{SubjectCanonical} \\ \vee \\ \text{SubjectRelative} \end{array}$

Object



Transitive Active Verb



Metagrammars as DCGs

- a metagrammar is a DCG
- terminals are tree descriptions

a metagrammar is the grammar of a grammar

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Model-Theoretic View

Set of tree descriptions \models Metagrammar
Lexical entry \models Set of tree descriptions

Operational View

- a metagrammar is a logic program
- its execution accumulates tree descriptions
- which are then processed by a solver
- resulting in the production of a lexical entry
- backtrack to obtain the rest of the lexicon

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1 Motivations

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Type definitions

```
type CAT={n,v,p}
```

```
type PERS=[1..3]
```

```
type FLEX=[num:NUMBER, gen:GENDER, pers:PERS]
```

Property definitions

```
property extraction : bool
```

Property definitions

```
property extraction : bool {extra = +}
```

```
extra ≡ extraction = +
```

Feature definitions

```
feature num : NUMBER
```

Class definitions

```
class C1
import C2[]
export ?X ?Y
declare ?X ?Y !Z
{
    Statement
}
```

Restricted imports

```
class C1
import C2[] as [?X1, ..., ?Xn]
export ?X ?Y
declare ?X ?Y !Z
{
    Statement
}
```

Renamings

```
class C1
import C2[] as [?X1=?Y1, ..., ?Xn]
export ?X=?U ?Y
declare ?X ?Y !Z
{
    Statement
}
```

Statements

$$\begin{aligned} S ::= \quad & S_1 ; S_2 \\ | \quad & S_1 \mid S_2 \\ | \quad & E_1 = E_2 \\ | \quad & \textit{Class}[E_1, \dots, E_n] \\ | \quad & \langle \textit{Dim} \rangle \{ \dots \} \\ | \quad & S *= [f = E, \dots] \end{aligned}$$

Expressions

$$\begin{array}{l} E ::= ?X \mid !X \\ \quad \mid Atom \mid Int \mid String \\ \quad \mid @\{Atom, \dots, Atom\} \\ \quad \mid ?X = [f_1 = E_1, \dots, f_n = E_n] \\ \quad \mid E_1.E_2 \\ \quad \mid E(E_1, \dots, E_n) \\ \quad \mid E_1 | E_2 \\ \quad \mid (E) \end{array}$$

Dimension-specific description languages

$\langle \text{Dim} \rangle \{ \dots \}$

- tree descriptions
- hole semantics
- AVM, (semi-)lattices

There are currently 3 built-in dimensions: `<syn>` `<sem>` `<dyn>`

Tree Description Language

$\langle \text{syn} \rangle \{ \text{Syn} \}$

$\text{Syn} ::= \text{Syn}; \text{Syn}$
| $\text{Syn} | \text{Syn}$
| **node** $?X \ (p = E, \dots) \ [f = E, \dots]$
| $E \rightarrow E \mid E \rightarrow^+ E \mid E \rightarrow^* E$
| $E \gg E \mid E \gg^+ E \mid E \gg^* E$
| $E = E$
| $E < E_1, \dots, E_n >$

Alternative Syntax

■ dominance

```
node {         node }  
node { ...+ node }  
node { ...   node }
```

■ precedence

```
node          node  
node , , , + node  
node , , ,   node
```

Flat Semantics Description Language

$\langle \text{sem} \rangle \{ \text{Sem} \}$

$\text{Sem} ::= \text{Sem}; \text{Sem}$
| $\text{Sem} | \text{Sem}$
| $\text{Var} : E(E, \dots, E)$
| $\text{Var} : \sim E(E, \dots, E)$
| $\text{Var} << \text{Var}$

Mutual Exclusion Sets

mutex SUBJ-INV

mutex SUBJ-INV += CanonicalObject

mutex SUBJ-INV += InvertedNominalSubject

2 classes in the same mutex cannot both be used in the same derivation

Principles